

General position problems in strong and lexicographic products of graphs

Pakanun Dokyeesun¹ · Sandi Klavžar^{1,2,3} · Dorota Kuziak⁴ · Jing Tian^{1,5}

Received: 28 August 2024 / Revised: 24 October 2025 / Accepted: 16 November 2025 © The Author(s) 2025

Abstract

Outer, dual, and total general position sets are studied on strong and lexicographic products of graphs. Sharp lower and upper bounds are proved for the outer and the dual general position number of strong products and several exact values are obtained. For the lexicographic product, the outer general position number is determined in all the cases, and the dual general position number is established in many cases. The total general position number is determined for both products. Along the way some results on outer general position sets are also derived.

Keywords Outer general position · Dual general position · Total general position · Strong product · Lexicographic product

Mathematics Subject Classification 05C12 · 05C69 · 05C76

1 Introduction

Let G = (V(G), E(G)) be a graph and $X \subseteq V(G)$. Vertices $u, v \in V(G)$ are X-positionable if for any shortest u, v-path P we have $V(P) \cap X \subseteq \{u, v\}$. Clearly, adjacent vertices are X-positionable no matter whether they belong to X. Setting $\overline{X} = V(G) \setminus X$, the set X is

• a general position set, if every $u, v \in X$ are X-positionable;

Sandi Klavžar sandi.klavzar@fmf.uni-lj.si

Pakanun Dokyeesun papakanun@gmail.com

Dorota Kuziak dorota.kuziak@uca.es

Jing Tian jingtian526@126.com

Published online: 15 December 2025

- Institute of Mathematics, Physics and Mechanics, Ljubljana, Slovenia
- Faculty of Mathematics and Physics, University of Ljubljana, Ljubljana, Slovenia
- Faculty of Natural Sciences and Mathematics, University of Maribor, Maribor, Slovenia
- Departamento de Estadística e Investigación Operativa, Universidad de Cádiz, Algeciras, Spain
- School of Science, Zhejiang University of Science and Technology, Hangzhou 310023, Zhejiang, People's Republic of China



97 Page 2 of 13 P. Dokyeesun et al.

- a total general position set, if every $u, v \in V(G)$ are X-positionable;
- an outer general position set, if every $u, v \in X$ are X-positionable, and every $u \in X$, $v \in \overline{X}$ are X-positionable; and
- a dual general position set, if every $u, v \in X$ are X-positionable, and every $u, v \in \overline{X}$ are X-positionable.

The cardinality of a largest general position set, a largest total general position set, a largest outer general position set, and a largest dual general position set of a graph G will be respectively denoted by gp(G), $gp_t(G)$, $gp_o(G)$, and $gp_d(G)$. These graph invariants will be respectively called the *general position number*, the *total general position number*, the *outer general position number*, and the *dual general position number* of G. Moreover, for any $\tau(G) \in \{gp(G), gp_t(G), gp_o(G), gp_d(G)\}$, by a τ -set we mean a set of cardinality $\tau(G)$ having the respective property.

The above described variety of general position sets was introduced in Tian and Klavžar (2025), while (standard) general position sets have already been investigated a lot after their independent introduction in Chandran and Parthasarathy (2016) and Manuel and Klavžar (2018). In particular, a lot of attention has been given to general position sets of graph operations (Ghorbani et al. 2021; Iršič et al. 2024; Klavžar and Rus 2021; Korže and Vesel 2023; Tian and Xu 2021; Tian et al. 2021). Among other contributions on the topic we highlight the papers Patkós (2020), Thomas et al. (2024) and Yao et al. (2022). With respect to graph operations, general position sets were also investigated in Klavžar and Yero (2019) on strong products. Among other matters covered in Klavžar and Yero (2019) we point to the problem whether the general position number of the strong product of two graphs equals the product of the general position numbers of the factors, which is one of the main outstanding problems in the area.

In this paper we focus on the recently introduced three general position sets of strong and lexicographic products. In the next section we give necessary definitions and call up the known results that we need later. In Sect. 3 we prove some results on outer general position sets that are on one hand interesting in their own right and are on the other hand applied later on. The strong product is the topic of Sect. 4. We prove general lower and upper bounds on the outer and the dual general position number, demonstrate their sharpness, and give exact values in some specific cases. Then, in Sect. 5, we turn our attention to lexicographic products. We first determine the outer general position number in all the cases, and complete the paper by establishing the dual general position number in many cases. Along the way we also determine the total general position number of both products. We conclude the article with some open problems that we find interesting for further research.

2 Preliminaries

Unless stated otherwise, graphs G = (V(G), E(G)) considered in this paper are simple and connected. For a positive integer k, we use [k] to represent the set $\{1, \ldots, k\}$. For a vertex u of G, $N_G(u)$ denotes the set of neighbors of u in G and $N_G[u] = N_G(u) \cup \{u\}$. Vertices u and v of G are true twins if $N_G[u] = N_G[v]$. True twins are thus adjacent vertices such that for each vertex $w \neq u$, v we have $uw \in E(G)$ if and only if $vw \in E(G)$. The degree of a vertex u is $d_G(u) = |N_G(u)|$ and the largest degree in the graph G is the maximum degree $\Delta(G)$. A vertex of G is universal if it is adjacent to all the other vertices of G. A vertex u of G is simplicial if $N_G(u)$ induces a complete subgraph. The set of simplicial vertices of G will be denoted by S(G) and the cardinality of S(G) by S(G). The order, the number of leaves, and



the complement of G are respectively denoted by n(G), $n_1(G)$, and \overline{G} . Moreover, $\omega(G)$ and $\alpha(G)$ stand for the clique number and the independence number of G.

The distance $d_G(u, v)$ between vertices u and v of G is the number of edges on a shortest u, v-path. The diameter of G is the maximum distance between pairs of vertices of G and is denoted by diam(G). Let H be a subgraph of G. Then H is isometric if for each pair of vertices $u, v \in V(H)$ we have $d_H(u, v) = d_G(u, v)$, and H is convex if for any vertices $u, v \in V(H)$, any shortest u, v-path in G lies completely in H. By abuse of language we will also say that a set $X \subseteq V(G)$ is convex if X induces a convex subgraph.

The vertex u of G is maximally distant from a vertex $v \in V(G)$ if every $w \in N_G(u)$ satisfies $d_G(v, w) \le d_G(u, v)$. If u is maximally distant from v, and v is maximally distant from u, then u and v are mutually maximally distant, MMD for short. Note that true twins are MMD. The set of MMD vertices of G is called the boundary of G and denoted by $\partial(G)$, see Brešar et al. (2008), Cáceres et al. (2005) and Kuziak et al. (2018), that is,

$$\partial(G) = \{u \in V(G) : u \text{ is an MMD vertex}\}.$$

We further set $b(G) = |\partial(G)|$. The strong resolving graph G_{SR} of G has V(G) as the vertex set, two vertices being adjacent in G_{SR} if they are MMD in G. The notion of the strong resolving graph was introduced in Oellermann and Peters-Fransen (2007) while investigating the strong metric dimension of graphs. Note that there could be vertices in a graph G which are not MMD with any other vertex in G (they do not belong to $\partial(G)$). Hence such vertices are isolated in G_{SR} . In this sense, from now on we denote by G'_{SR} the graph obtained from G_{SR} by removing all (possibly zero) of its isolated vertices. Note that $V(G'_{SR}) = \partial(G)$.

We next collect characterizations of the three newly introduced general position sets as respectively proved in Tian and Klavžar (2025, Theorems 2.1, 2.3, 3.1).

Theorem 2.1 If G is a connected graph and $X \subseteq V(G)$, then the following hold.

- (i) *X* is a total general position set of *G* if and only if $X \subseteq S(G)$. Moreover, $\operatorname{gp}_{\tau}(G) = s(G)$.
- (ii) If $|X| \ge 2$, then X is an outer general position set of G if and only if each pair of vertices from X is mutually maximally distant. Moreover, $\operatorname{gp}_0(G) = \omega(G_{SR})$.
- (iii) If X is a general position set of G, then X is a dual general position set if and only if G X is convex.

Notice that by the definition of G'_{SR} it follows that $\omega(G'_{SR}) = \omega(G_{SR})$. Moreover, by Theorem 2.1(ii), we have that for any connected graph G,

$$gp_0(G) = \omega(G'_{SR}). \tag{1}$$

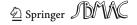
3 On outer general position sets

In this section we prove some results on outer general position sets that are interesting in their own right and that we will need later in the paper. The first of them holds for all types of general position sets.

Lemma 3.1 Let H be an isometric subgraph of a graph G. If X is a (dual, outer, total) general position set of G, then $X \cap V(H)$ is a (dual, outer, total) general position set of H.

Proof Set $X_H = X \cap V(H)$.

Assume that X is a general position set of G and suppose on the contrary that X_H is not a general position set of H. Then there exist vertices h, h', h'' of X_H such that h'' lies on a



97 Page 4 of 13 P. Dokyeesun et al.

shortest h, h'-path P in H. Since H is isometric in G, the path P is also a shortest path in G, but then X is not a general position set of G. This contradiction shows that X_H is a general position set of H.

Assume now that X is a dual general position set of G. We already know that X_H is a general position set of H, hence in view of Theorem 2.1(iii) we claim that $V(H)\backslash X_H$ is convex in H. Assume on the contrary that this is not the case. Then there exist two vertices $u, v \in V(H)\backslash X_H$ and a shortest u, v-path Q in H which contains a vertex of X_H . As H is isometric, Q is also a shortest path in G. But since $u, v \in V(G)\backslash X$, this implies that $V(G)\backslash X$ is not convex in G, a contradiction with Theorem 2.1(iii). It follows that X_H is a dual general position set of H.

Assume next that X is an outer general position set of G. Then X_H is a general position set of H and similarly as above we infer that also no shortest path between a vertex from X_H and a vertex from $V(H)\backslash X_H$ contains a vertex of X_H . It follows that X_H is an outer general position set of H.

Finally, if X is a total general position set of G, then the above arguments imply that X_H is a total general position set of H.

The case of Lemma 3.1 for general position sets was earlier implicitly proved within the proof of Manuel and Klavžar (2018, Theorem 3.1).

If $k \ge 1$, then a set $X \subseteq V(G)$ is *k-independent* if $d_G(u, v) > k$ for each pair of vertices $u, v \in X$. The *k-independence number* of G is the cardinality of a largest *k*-independent set in G and denoted by $\alpha_k(G)$. Note that $\alpha_1(G) = \alpha(G)$.

Proposition 3.2 If G is a graph with diam $(G) = k \ge 2$, then $gp_0(G) \ge \alpha_{k-1}(G)$.

Proof Let X be a (k-1)-independent set of G with $|X| = \alpha_{k-1}(G)$. Let $u, v \in X$, where $u \neq v$. As $\operatorname{diam}(G) = k$ and $d_G(u, v) > k-1$, we necessarily have $d_G(u, v) = k$. Using again the fact that $\operatorname{diam}(G) = k$ we see that u and v are MMD. It follows that X induces a clique of G_{SR} , so that $\omega(G_{SR}) \geq |X| = \alpha_{k-1}(G)$. Theorem 2.1(ii) completes the argument.

The bound of Proposition 3.2 is sharp for all $k \ge 2$. To see it, we give two constructions depending on the parity of k.

First, let $k \ge 2$ be even and let $r = \frac{k-2}{2}$. Then the graph $K_{1,s}^r$, $s \ge 2$, is obtained from $K_{1,s}$ by subdividing all edges r times. Notice that $\operatorname{diam}(K_{1,s}^r) = k = 2r + 2$ and all leaves of this graph form the largest (2r+1)-independence set. Hence $\alpha_{k-1}(K_{1,s}^r) = n_1(K_{1,s}^r)$. Moreover, by Theorem 2.1(ii) we have that $\operatorname{gp}_0(K_{1,s}^r) = \omega((K_{1,s}^r)_{\operatorname{SR}}) = n_1(K_{1,s}^r) = \alpha_{k-1}(K_{1,s}^r)$.

For $k \ge 3$ odd, let $t = \frac{k-1}{2}$. Then the graph K_n^t , $n \ge 2$, is obtained by attaching a path P_t to every vertex of K_n . By construction, diam $(K_n^t) = k = 2t + 1$ and $\alpha_{k-1}(K_n^t) = n_1(K_n^t)$. Using Theorem 2.1(ii) again we can conclude that $gp_0(K_n^t) = \omega((K_n^t)_{SR}) = n_1(K_n^t) = \alpha_{k-1}(K_n^t)$.

By G_{tt}^- we represent the graph obtained from G by removing the edge between each pair of true twins. In particular, if G has no true twins, then $G_{\operatorname{tt}}^- \cong G$, and $(K_n)_{\operatorname{tt}}^-$ is the edgeless graph of order n.

Proposition 3.3 *If* G *is a graph with* $\operatorname{diam}(G) = 2$, then $\operatorname{gp}_0(G) = \alpha(G_{\operatorname{tt}}^-)$. In particular, if G has no true twins, then $\operatorname{gp}_0(G) = \alpha(G)$.

Proof We have already observed that every pair of true twins are MMD in a graph. Since diam(G) = 2, vertices u and v of G are MMD if and only if either $d_G(u, v) = 2$, or u and v are true twins. Then for every $u, v \in V(G)$, vertices u, v are MMD in G if and only if they



nt.

are MMD in G_{tt}^- which means that $G_{SR} \cong (G_{tt}^-)_{SR}$. By Theorem 2.1(ii) and diam(G) = 2, we have $gp_o(G) = \omega(G_{SR}) = \omega((G_{tt}^-)_{SR}) = \omega(G_{tt}^-) = \alpha(G_{tt}^-)$.

Note that any two non-adjacent vertices are MMD as $\operatorname{diam}(G) = 2$. If G has no true twins, we have $\omega(G_{SR}) = \omega(\overline{G}) = \alpha(G)$ and hence $\operatorname{gp}_0(G) = \alpha(G)$.

When G has no true twins, Proposition 3.3 demonstrates that the bound established in Proposition 3.2 is sharp for k = 2.

A parallel result to Proposition 3.3 for the case of general position sets was proved in Klavžar and Yero (2019, Proposition 2.4), where it is established that if G has no true twins and diam(G) = 2, then $gp(G) = \omega(G_{SR})$ if and only if $gp(G) = \alpha(G)$.

4 Strong products

The *strong product* of two graphs G and H is the graph $G \boxtimes H$ such that $V(G \boxtimes H) = V(G) \times V(H)$, and two vertices $(g,h), (g',h') \in V(G \boxtimes H)$ are adjacent if either g = g' and $hh' \in E(H)$, or $gg' \in E(G)$ and h = h', or $gg' \in E(G)$ and $hh' \in E(H)$. For a vertex $h \in V(H)$, set $G^h = \{(g,h) \in V(G \boxtimes H) : g \in V(G)\}$. The set G^h is called a G-layer of $G \boxtimes H$ and induces a subgraph of $G \boxtimes H$ isomorphic to G. For $g \in V(G)$, G is defined as $G = \{(g,h) \in V(G \boxtimes H) : h \in V(H)\}$. If $G \subseteq V(G \boxtimes H)$, the projection $G \subseteq V(G)$ of $G \subseteq V(G)$ to $G \subseteq V(G)$ is the set $G \subseteq V(G)$ in $G \subseteq V(G)$. The projection $G \subseteq V(G)$ is defined analogously.

The following formulas are well-known, cf. Hammack et al. (2011):

$$d_{G \boxtimes H}((g, h), (g', h')) = \max\{d_G(g, g'), d_H(h, h')\}$$
 and (2)

$$N_{G \boxtimes H}[(g,h)] = N_G[g] \times N_H[h]. \tag{3}$$

In Klavžar and Yero (2019), several exact results and bounds were proved on the general position number of strong products. In addition, the problem was posed whether $gp(G \boxtimes H) = gp(G)gp(H)$ holds for any graphs G and H (Klavžar and Yero 2019, Problem 4.8). In this section we consider the other three versions of general position sets on strong products.

The next result follows from (3) and can also be deduced from the proof of Cáceres et al. (2010, Proposition 4).

Lemma 4.1 If G and H are graphs, then $S(G \boxtimes H) = S(G) \times S(H)$.

Lemma 4.1 together with Theorem 2.1(i) yield:

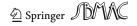
Corollary 4.2 If G and H are graphs, then $gp_t(G \boxtimes H) = s(G)s(H)$.

4.1 Outer general position sets

To prove bounds on the outer general position number of strong products, we recall the following result that clarifies MMD vertices of strong products.

Lemma 4.3 (Oellermann and Peters-Fransen 2007, Lemma 2.6) *Let G and H be two connected graphs. Let g, g' be two vertices of G and h, h' two vertices of H. Then* (g, h) *and* (g', h') *are MMD in G* \boxtimes *H if and only if one the following conditions holds:*

- (i) g, g' are MMD in G and h, h' are MMD in H;
- (ii) g, g' are MMD in G and h = h';



97 Page 6 of 13 P. Dokyeesun et al.

- (iii) h, h' are MMD in H and g = g';
- (iv) g, g' are MMD in G and $d_G(g, g') > d_H(h, h')$;
- (v) h, h' are MMD in H and $d_G(g, g') < d_H(h, h')$.

Recall that the boundary $\partial(G)$ of G contains its MMD vertices and that $b(G) = |\partial(G)|$. The bounds on the outer general position number of strong products now read as follows.

Theorem 4.4 If G and H are connected graphs of order at least 2, then

$$\operatorname{gp}_{0}(G)\operatorname{gp}_{0}(H) \leq \operatorname{gp}_{0}(G \boxtimes H) \leq b(G)b(H)$$
.

Moreover, if G and H are block graphs, then the bounds coincide.

Proof To prove the lower bound, let X be a gp_o -set of G and Y be a gp_o -set of H. We will prove that $X \times Y$ is an outer general position set of $G \boxtimes H$. Consider any two vertices (g,h), (g',h') from $X \times Y$. Assume first that $g \neq g'$ and $h \neq h'$. Since X and Y are gp_o -sets of G and H, Theorem 2.1(ii) implies that g,g' are MMD in G and G and G and G are MMD in G and G and G are MMD in G are MMD

To prove the upper bound, let *S* be a gp_o-set of $G \boxtimes H$ and consider the following cases. **Case 1**: $p_G(S) \subseteq \partial(G)$ and $p_H(S) \subseteq \partial(H)$.

In this case $\operatorname{gp}_{o}(G \boxtimes H) = |S| \leq b(G)b(H)$.

Case 2: S contains a vertex (g, h) such that $h \notin \partial(H)$.

Let (g',h') be an arbitrary vertex from S, where $g' \neq g$. Since S is an outer general position set of $G \boxtimes H$, Theorem 2.1(ii) implies that (g,h) and (g',h') are MMD in $G \boxtimes H$. Since $h \notin \partial(H)$, Lemma 4.3 yields that g and g' are MMD in G. We thus see that every vertex of $p_G(S) \setminus \{g\}$ is adjacent to g in the strong resolving graph G_{SR} and hence $|p_G(S) \setminus \{g\}| \leq \Delta(G_{SR})$. Using Lemma 4.3 once more we see that each H-layer g'H has at most $gp_O(H)$ vertices belonging to S, that is, $|V(g'H) \cap S| \leq gp_O(H)$. Moreover, since $h \notin \partial(H)$, we have $|V(gH) \cap S| = 1$. Indeed, for otherwise a vertex $(g,h') \in S$, where $h' \neq h$, would imply that $h \in \partial(H)$ by Lemma 4.3(iii). Then $|S| \leq \Delta(G_{SR})gp_O(H) + 1$. By definitions of boundary sets and strong resolving graphs, we have $\Delta(G_{SR}) \leq b(G) - 1$ and $\omega(H_{SR}) \leq b(H)$. Together with Theorem 2.1(ii) we get

$$\begin{split} |S| & \leq \Delta(G_{\text{SR}}) \text{gp}_{\text{o}}(H) + 1 \\ & \leq (b(G) - 1)\omega(H_{\text{SR}}) + 1 \\ & \leq (b(G) - 1)b(H) + 1 \\ & < b(G)b(H). \end{split}$$

Case 3: *S* contains a vertex (g, h) such that $g \notin \partial(G)$.

By the commutativity of the strong product we can use parallel arguments as in Case 2 to derive the conclusion $|S| \le \Delta(H_{SR}) gp_o(G) + 1 < b(G)b(H)$.

Case 4: S contains a vertex (g, h) such that $g \notin \partial(G)$ and $h \notin \partial(H)$.

Lemma 4.3 implies that one of g and h must actually respectively lie in $\partial(G)$ and $\partial(H)$, hence this case is not possible.

We have exhausted all the possibilities for S, hence the claimed upper bound is proved.



Assume now that G and H are block graphs. Then we know from Tian and Klavžar (2025) that $gp_0(G) = s(G)$ and $gp_0(H) = s(H)$. Since in a block graph every two simplicial vertices are MMD, we get b(G) = s(G) and b(H) = s(H). Hence the bounds of the theorem coincide if G and H are block graphs.

The bounds of Theorem 4.4 are thus sharp on block graphs. On the other hand, recall from Sonnemann and Krafft (1974) that $\alpha(C_5 \boxtimes C_5) = 5$. By (2) we get diam($C_5 \boxtimes C_5$) = 2, hence Proposition 3.3 yields $\operatorname{gp}_0(C_5 \boxtimes C_5) = 5$, the value which lies strictly between $\operatorname{gp}_0(C_5)\operatorname{gp}_0(C_5) = 4$ and $b(C_5)b(C_5) = 25$.

Denoting the strong product of k copies of G by $G^{k,\boxtimes}$, the last result of this subsection reads as follows.

Proposition 4.5 *If* G *is a diameter two graph without true twins and* $k \ge 1$, *then* $gp_0(G^{k, \boxtimes}) = \alpha(G^{k, \boxtimes})$.

Proof If k=1, the assertion is just Proposition 3.3. Since G has no true twins, we can deduce by applying (Barragán-Ramírez and Rodríguez-Velázquez 2016, Lemma 9) that $G \boxtimes G$ also has no true twins. By induction on k, the power $G^{k,\boxtimes}$ also has no true twins for each $k \ge 2$. Moreover, by (2) and induction we also have $\operatorname{diam}(G^{k,\boxtimes}) = 2$. Applying Proposition 3.3 again we reach the conclusion.

4.2 Dual general position sets

For the dual general position number of strong products we have the following general bounds.

Theorem 4.6 *If G and H are connected graphs, then*

$$s(G)s(H) \le \operatorname{gp}_{\operatorname{d}}(G \boxtimes H) \le \min\{s(G)n(H) + s(H)n(G) - s(G)s(H), \\ n(G)\operatorname{gp}_{\operatorname{d}}(H), n(H)\operatorname{gp}_{\operatorname{d}}(G)\}.$$

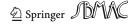
Proof In Tian and Klavžar (2025, Corollary 3.2) it was proved that if G is a graph and $X \subseteq S(G)$, then X is a dual general position set of G. Hence $S(G \boxtimes H)$ is a dual general position set of $G \boxtimes H$. Lemma 4.1 then yields the lower bound.

Lemma 3.1 applied to dual general position sets implies that $\operatorname{gp}_{\operatorname{d}}(G \boxtimes H) \leq n(G)\operatorname{gp}_{\operatorname{d}}(H)$ and $\operatorname{gp}_{\operatorname{d}}(G \boxtimes H) \leq n(H)\operatorname{gp}_{\operatorname{d}}(G)$.

To prove that $\operatorname{gp}_{\operatorname{d}}(G\boxtimes H)\leq s(G)n(H)+s(H)n(G)-s(G)s(H)$, suppose on the contrary that there exists a dual general position set R of $G\boxtimes H$ such that |R|>s(G)n(H)+s(H)n(G)-s(G)s(H). Then there exists a vertex $(x,y)\in R$ such that $x\notin S(G)$ and $y\notin S(H)$, that is, x and y are not simplicial vertices of G and G, respectively. Let $g',g''\in N_G(x)$, and let $g',g''\in N_G(x)$, where $g',g'''\notin S(G)$ and $g',g'''\notin S(G)$. Then the subgraph induced by g',g',g',g'' is isomorphic to $g',g''\in S(G)$ with the universal vertex g',g',g'.

Since R is a dual general position set of $G \boxtimes H$ and because $(x, y) \in R$, Theorem 2.1(iii) implies that one of the vertices (x', y) and (x'', y) must be in R, we may without loss of generality assume that $(x', y) \in R$. Using again that R is a general position set of $G \boxtimes H$, we get $(x'', y') \notin R$ and $(x'', y'') \notin R$. Since R is a dual general position set of G, by Theorem 2.1(iii), G - R is convex. But then the shortest (x'', y'), (x'', y'')-path contains the vertex (x, y), a contradiction with the assumption that R is a dual general position.

To demonstrate that the three upper bounds of Theorem 4.6 are pairwise incomparable, consider the following example. Let C_n^+ be the graph obtained from C_n by attaching a new



97 Page 8 of 13 P. Dokyeesun et al.

vertex to an arbitrary vertex of C_n . Now we consider $C_{2k+1}^+ \boxtimes C_{2\ell+1}^+$, $k, \ell \ge 2$, and compute the values of the upper bounds of Theorem 4.6. Since $s(C_{2k+1}^+) = s(C_{2\ell+1}^+) = 1$, the first bound is

$$s(C_{2k+1}^+)n(C_{2\ell+1}^+) + s(C_{2\ell+1}^+)n(C_{2k+1}^+) - s(C_{2k+1}^+)s(C_{2\ell+1}^+)$$

$$= (2\ell+2) + (2k+2) - 1$$

$$= 2\ell + 2k + 3.$$

For the other two bounds, we first infer that $\operatorname{gp}_{\operatorname{d}}(C_{2k+1}^+)=3$ and $\operatorname{gp}_{\operatorname{d}}(C_{2\ell+1}^+)=3$. Assume without loss of generality that $k<\ell$. Then we have

$$\min\{\operatorname{gp}_{\operatorname{d}}(C_{2k+1}^+)n(C_{2\ell+1}^+),\operatorname{gp}_{\operatorname{d}}(C_{2\ell+1}^+)n(C_{2k+1}^+)\}=\min\{6\ell+6,6k+6\}=6k+6.$$

Theorem 4.6 thus gives

$$\operatorname{gp}_{\operatorname{d}}(C_{2k+1}^{+} \boxtimes C_{2\ell+1}^{+}) \leq \begin{cases} 2\ell + 2k + 3; & k \leq \ell \leq 2k + \frac{3}{2}, \\ 6k + 6; & \ell > 2k + \frac{3}{2}. \end{cases}$$

We conclude this section with the following exact result on the dual general position number of strong products.

Theorem 4.7 *If* H *is a connected graph, then* $gp_d(K_n \boxtimes H) = n \cdot gp_d(H)$.

Proof By Theorem 4.6, $\operatorname{gp}_{\operatorname{d}}(K_n \boxtimes H) \leq n \cdot \operatorname{gp}_{\operatorname{d}}(H)$.

To prove the lower bound, let $V(K_n) = [n]$ and let X be a gp_d -set of H. We claim that $X' = [n] \times X$ is a dual general position set of $K_n \boxtimes H$. We first need to show that X' is a general position set of $K_n \boxtimes H$. For the sake of contradiction suppose this is not the case and let (i,h), (i',h'), (i'',h'') be three vertices of X' such that (i'',h'') lies on a shortest (i,h), (i',h')-path P in $K_n \boxtimes H$. Then $p_H(P)$ induces a shortest h,h'-path in H which contains h'', a contradiction with the assumption that X is a general position set of H.

We second need to verify that $V(K_n \boxtimes H) \setminus X'$ is convex in $K_n \boxtimes H$. For this sake consider arbitrary two vertices (i, h) and (i', h') from $V(K_n \boxtimes H) \setminus X'$. Suppose that there exists a shortest (i, h), (i', h')-path P containing a vertex from X', say (j, k), such that $k \in X$. Then it follows that there exists a shortest h, h'-path containing the vertex k in H. Since $V(H) \setminus X$ is convex, $d_H(h, h') < d_H(h, k) + d_H(k, h')$. By (2), we have

$$d_{K_n \boxtimes H}((i, h), (i', h')) = d_H(h, h'),$$

$$d_{K_n \boxtimes H}((i, h), (j, k)) = d_H(h, k),$$

$$d_{K_n \boxtimes H}((j, k), (i', h')) = d_H(k, h').$$

Then

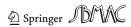
$$d_{K_n \boxtimes H}((i,h),(i',h')) < d_{K_n \boxtimes H}((i,h),(j,k)) + d_{K_n \boxtimes H}((j,k),(i',h')),$$

which implies that P is not a shortest (i, h), (i', h')-path in $K_n \boxtimes H$. This contradiction yields that $V(K_n \boxtimes H) \setminus X'$ is convex. By Theorem 2.1(iii) we can conclude that X' is a dual general position set of $K_n \boxtimes H$ and hence $\operatorname{gp}_d(K_n \boxtimes H) \ge |X'| = n \cdot \operatorname{gp}_d(H)$, and we are done. \square

Let H be an arbitrary graph with $gp_d(H) = s(H)$, for instance, an arbitrary block graph. Then Theorem 4.7 yields

$$\operatorname{gp}_{\operatorname{d}}(K_n \boxtimes H) = n \cdot \operatorname{gp}_{\operatorname{d}}(H) = s(K_n)s(H),$$

which demonstrates sharpness of the lower bound of Theorem 4.6.



5 Lexicographic products

The *lexicographic product* $G \circ H$ of two graphs G and H has $V(G \circ H) = V(G) \times V(H)$. Two vertices $(g,h), (g',h') \in V(G \circ H)$ are adjacent if either g=g' and $hh' \in E(H)$, or $gg' \in E(G)$. The projections $p_G(X)$ and $p_H(X)$ of $X \subseteq V(G \circ H)$ on G and H are defined just as for the strong product, and G-layers and H-layers are also defined analogously.

The following result is parallel to Lemma 4.1. Since we did not find it in the literature, its proof is provided.

Lemma 5.1 If G and H are connected graphs of order at least 2, then

$$S(G \circ H) = \begin{cases} S(G) \times V(H); & H \text{ is complete,} \\ \emptyset; & \text{otherwise.} \end{cases}$$

Proof We first assume that H is a non-complete graph and suppose on the contrary that $S(G \circ H) \neq \emptyset$. Let $(x, y) \in S(G \circ H)$. As G is a connected graph with at least two vertices, there exists $x' \in V(G)$ such that $xx' \in E(G)$. Since H is non-complete, there exist two non-adjacent vertices $y', y'' \in V(H)$. Then $(x', y'), (x', y'') \in N_{G \circ H}((x, y))$ but they are not adjacent, a contradiction.

Assume second that H is complete. Consider a vertex $(x, y) \in S(G) \times V(H)$ and let $(x', y'), (x'', y'') \in N_{G \circ H}((x, y))$. If x' = x'', then (x', y') and (x'', y'') are adjacent since H is complete. And if $x' \neq x''$, then (x', y') and (x'', y'') are adjacent because $x \in S(G)$ and hence $y'y'' \in E(H)$. This proves that $S(G) \times V(H) \subseteq S(G \circ H)$. To see that also $S(G \circ H) \subseteq S(G) \times V(H)$ holds, we can use the argument from the first paragraph of the proof.

Lemma 5.1 together with Theorem 2.1(i) yield:

Corollary 5.2 *If G and H are graphs of order at least* 2, *then*

$$gp_{t}(G \circ H) = \begin{cases} s(G)n(H); & H \text{ is complete,} \\ 0; & \text{otherwise.} \end{cases}$$

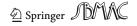
5.1 Outer general position sets

In this subsection we determine the outer general position number of lexicographic products. For this sake, we first need to recall some concepts and results presented in Kuziak et al. (2016). Moreover, the fact $\omega(G \circ H) = \omega(G)\omega(H)$, see Došlić et al. (2013, Theorem 3), is used several times in this subsection.

Given a graph G, by $G_{\overline{2}}$ it is denoted the graph with vertex set $V(G_{\overline{2}}) = V(G)$ such that two vertices u, v are adjacent in $G_{\overline{2}}$ if either $d_G(u, v) \geq 2$ or u, v are true twins. Further, by G' we denote the graph obtained from G by removing all its isolated vertices. (Note that this notation is consistent with G'_{SR} used earlier for strong resolving graphs.) The next result is a direct consequence of Kuziak et al. (2016, Remark 6), where G + H denotes the join of graphs G and H.

Proposition 5.3 *Let G be a connected graph of order at least* 2.

- (i) If G has no universal vertex, then $\omega(G_{\overline{2}}) = \omega((K_1 + G)'_{SR})$.
- (ii) If diam $(G) \le 2$, then $\omega(G'_{\overline{2}}) = \omega(G'_{SR})$.



97 Page 10 of 13 P. Dokyeesun et al.

(iii) If G has no true twins, then $\omega(G_{\overline{2}}) = \alpha(G)$.

We next collect (Kuziak et al. 2016, Propositions 8, 11, 13, 16) into the following result. Before that some extra preparation is needed. The *TF-boundary* of a non-complete graph G is a set $\partial_{TF}(G) \subseteq \partial(G)$, where $x \in \partial_{TF}(G)$ whenever there exists $y \in \partial(G)$, such that x and y are MMD in G and $N_G[x] \neq N_G[y]$ (x, y are not true twins). Moreover, the *strong resolving TF-graph* of G is a graph G_{SRS} with vertex set $V(G'_{SRS}) = \partial_{TF}(G)$, where two vertices x, y are adjacent in G'_{SRS} if x and y are MMD in G and $N_G[x] \neq N_G[y]$. Finally, if F is a subgraph of G, then for a given vertex $v \in V(G)$, by $(F)_v$ we denote the subgraph of $G \circ H$ isomorphic to F which lies in the H-layer of $G \circ H$ corresponding to the vertex v of G.

Proposition 5.4 (Kuziak et al. 2016, Propositions 8, 11, 13, 16) *If G and H are connected graphs of order at least 2, then the following properties hold.*

(i) If G has no true twins and H is non-complete, then

$$(G \circ H)'_{\operatorname{SR}} \cong \left(G'_{\operatorname{SR}} \circ H_{\overline{2}}\right) \cup \bigcup_{v \in V(G) \setminus \partial(G)} (H'_{\overline{2}})_v.$$

(ii) If m > 2, then

$$(G \circ K_m)'_{\mathrm{SR}} \cong (G'_{\mathrm{SR}} \circ K_m) \cup \bigcup_{v \in V(G) \setminus \partial(G)} (K_m)_v.$$

(iii) If m > 2 and H has no universal vertex, then

$$(K_m \circ H)'_{\operatorname{SR}} \cong \bigcup_{v \in V(K_m)} (H_{\overline{2}})_v.$$

(iv) If G is non-complete and H has no universal vertex, then

$$(G \circ H)'_{\mathsf{SR}} \cong (G'_{\mathit{SRS}} \circ H_{\overline{2}}) \cup \bigcup_{v \in V(G) \setminus \partial_{\mathit{TF}}(G)} (H_{\overline{2}})_v.$$

With Proposition 5.4 in hand, we can deduce the subsequent results. In the first one, we avoid true twins in the first factor of the product, and complete graph in the second one.

Theorem 5.5 Let G be a connected graph without true twins of order at least 2 and let H be a non-complete graph.

- (i) If H has no universal vertex, then $gp_0(G \circ H) = gp_0(G)gp_0(K_1 + H)$.
- (ii) If $\operatorname{diam}(H) = 2$, then $\operatorname{gp}_0(G \circ H) = \operatorname{gp}_0(G)\operatorname{gp}_0(H)$.
- (iii) If H has no true twins, then $gp_0(G \circ H) = gp_0(G)\alpha(H)$.

Proof By (1) and Proposition 5.4(i), we have

$$\begin{split} \operatorname{gp}_{\operatorname{o}}(G \circ H) &= \omega((G \circ H)'_{\operatorname{SR}}) = \omega \left(\left(G'_{\operatorname{SR}} \circ H_{\overline{2}} \right) \cup \bigcup_{v \in V(G) \setminus \partial(G)} (H'_{\overline{2}})_v \right) \\ &= \omega \left(G'_{\operatorname{SR}} \circ H_{\overline{2}} \right) = \omega(G'_{\operatorname{SR}}) \omega(H_{\overline{2}}) \,. \end{split}$$

Combining Proposition 5.3 with (1) we get the conclusion.



Theorem 5.6 If G is a connected graph of order at least 2 and m > 2, then

$$gp_o(G \circ K_m) = m \cdot gp_o(G).$$

Proof By (1) and Proposition 5.4(ii), we have

$$gp_{o}(G \circ K_{m}) = \omega((G \circ K_{m})'_{SR}) = \omega \left(\left(G'_{SR} \circ K_{m} \right) \cup \bigcup_{v \in V(G) \setminus \partial(G)} (K_{m})_{v} \right)$$
$$= \omega \left(G'_{SR} \circ K_{m} \right) = \omega(G'_{SR}) \omega(K_{m}).$$

Using (1) once more, we get the formula.

We now supplement the above results with the following.

Theorem 5.7 Let H be a graph of order at least 2 without universal vertex.

- 1. If m > 2, then the following hold.
 - (i) If diam(H) = 2, then $gp_o(K_m \circ H) = gp_o(H)$.
 - (ii) If diam(H) > 2, then $gp_0(K_m \circ H) = gp_0(K_1 + H)$.
- 2. If G is a connected non-complete graph, then the following hold.

 - (i) If $\operatorname{diam}(H)=2$, then $\operatorname{gp}_0(G\circ H)=\omega(G'_{SRS})\operatorname{gp}_0(H)$. (ii) If $\operatorname{diam}(H)>2$, then $\operatorname{gp}_0(G\circ H)=\omega(G'_{SRS})\operatorname{gp}_0(K_1+H)$.

Proof 1. By (1) and Proposition 5.4(iii), we have

$$\operatorname{gp}_{0}(K_{m} \circ H) = \omega((K_{m} \circ H)'_{\operatorname{SR}}) = \omega\left(\bigcup_{v \in V(K_{m})} (H_{\overline{2}})_{v}\right) = \omega(H_{\overline{2}}).$$

Therefore, by Proposition 5.3 and again (1) we conclude the proof for this situation.

2. Using (1) and Proposition 5.4(iv), we get

$$\begin{split} \mathrm{gp}_{\mathrm{o}}(G \circ H) &= \omega((G \circ H)'_{\mathrm{SR}}) = \omega \left((G'_{SRS} \circ H_{\overline{2}}) \cup \bigcup_{v \in V(G) \setminus \partial_{TF}(G)} (H_{\overline{2}})_v \right) \\ &= \omega \left(G'_{SRS} \circ H_{\overline{2}} \right) = \omega(G'_{SRS}) \omega(H_{\overline{2}}) \,, \end{split}$$

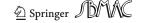
and then Proposition 5.3 and (1) yield the desired conclusion.

5.2 Dual general position sets

We finally investigate the dual general position number. For this sake we recall that in Anand et al. (2012, Theorem 2.1), non-complete convex sets of lexicographic products $G \circ H$ are characterized with three conditions. One of them asserts that $G \circ H$ can contain a non-complete convex subgraph only if H is a complete graph.

Theorem 5.8 *If G is a connected graph, then the following hold.*

- (i) If G and H are graphs without simplicial vertices, then $gp_d(G \circ H) = 0$.
- (ii) If $n \ge 1$, then $\operatorname{gp}_{\operatorname{d}}(G \circ K_n) = \operatorname{gp}_{\operatorname{d}}(G) \cdot n$.



97 Page 12 of 13 P. Dokyeesun et al.

Proof (i) Let X be a dual general position set of $G \circ H$ and let $(g,h) \in X$. Since $g \notin S(G)$, it has neighbors g' and g'' such that $g'g'' \notin E(G)$. Consider now the vertices (g',h) and (g'',h) and note that $d_{G \circ H}((g',h),(g'',h))=2$. Thus we cannot have $(g',h) \in X$ and $(g'',h) \in X$. Moreover, we also cannot have $(g',h) \notin X$ and $(g'',h) \notin X$ because $(G \circ H) \setminus X$ is convex. Hence, assume without loss of generality that $(g',h) \in X$ and $(g'',h) \notin X$. Let now (g'',h') be a vertex with $d_H(h,h')=2$. Such a vertex exists since h is not simplicial. Then $(g',h) \in X$ and $(g,h) \in X$ imply that $(g'',h') \notin X$. But then (g'',h')-(g,h)-(g'',h) is a shortest path with $(g,h) \in X$, a contradiction with the convexity of $(G \circ H) \setminus X$.

(ii) Since $G \circ K_n$ is isomorphic to $K_n \boxtimes G$, this follows by Theorem 4.7.

6 Open problems

This paper establishes that $\operatorname{gp}_{\mathsf{t}}(G \circ K_n) = \operatorname{gp}_{\mathsf{o}}(G \circ K_n) = \operatorname{gp}_{\mathsf{d}}(G \circ K_n)$ holds when G is a block graph. A natural direction for future research would be to characterize other graph classes for which this equality remains valid.

This study provides a comprehensive determination of the total and outer general position numbers for the lexicographic product $K_m \circ G$, see Corollary 5.2 and Theorem 5.7. It remains to determine the dual general position numbers for the lexicographic product $K_m \circ G$. The situation here seems more complicated, so we leave this as an open problem.

Acknowledgements This work was supported by the Slovenian Research Agency ARIS (research core funding P1-0297 and projects N1-0285, N1-0355, N1-0431). D. Kuziak has been partially supported by the Spanish Ministry of Science and Innovation through the Grant PID2023-146643NB-I00. D. Kuziak was visiting the University of Ljubljana supported by "Ministerio de Educación, Cultura y Deporte", Spain, under the "José Castillejo" program for young researchers (reference number: CAS22/00081).

Data availability Our manuscript has no associated data.

Declarations

Conflict of interest The authors declare that they have no known competing financial interests or personal relationships that could have appeared to influence the work reported in this paper.

Open Access This article is licensed under a Creative Commons Attribution 4.0 International License, which permits use, sharing, adaptation, distribution and reproduction in any medium or format, as long as you give appropriate credit to the original author(s) and the source, provide a link to the Creative Commons licence, and indicate if changes were made. The images or other third party material in this article are included in the article's Creative Commons licence, unless indicated otherwise in a credit line to the material. If material is not included in the article's Creative Commons licence and your intended use is not permitted by statutory regulation or exceeds the permitted use, you will need to obtain permission directly from the copyright holder. To view a copy of this licence, visit http://creativecommons.org/licenses/by/4.0/.

References

Anand BS, Changat M, Klavžar S, Peterin I (2012) Convex sets in lexicographic products of graphs. Graphs Combin 28:77–84

Barragán-Ramírez GA, Rodríguez-Velázquez JA (2016) The local metric dimension of strong product graphs. Graphs Combin 32:1263–1278

Brešar B, Klavžar S, Tepeh Horvat A (2008) On the geodetic number and related metric sets in Cartesian product graphs. Discrete Math 308:5555–5561



- Cáceres J, Puertas ML, Hernando C, Mora M, Pelayo IM, Seara C (2005) Searching for geodetic boundary vertex sets. Electron Notes Discrete Math 19:25–31
- Cáceres J, Hernando C, Mora M, Pelayo IM, Puertas ML (2010) On the geodetic and the hull numbers in strong product graphs. Comput Math Appl 60:3020–3031
- Chandran S.V, Parthasarathy GJ (2016) The geodesic irredundant sets in graphs. Int J Math Combin 4:135–143 Došlić T, Ghorbani M, Hosseinzadeh MA (2013) The relationships between Wiener index, stability number and clique number of composite graphs. Bull Malays Math Sci Soc 36:165–172
- Ghorbani M, Klavžar S, Maimani HR, Momeni M, Rahimi-Mahid F, Rus G (2021) The general position problem on Kneser graphs and on some graph operations. Discuss Math Graph Theory 41:1199–1213
- Hammack R, Imrich W, Klavžar S (2011) Handbook of product graphs, 2nd edn. CRC Press, Boca Raton Iršič V, Klavžar S, Rus G, Tuite J (2024) General position polynomials. Results Math 79:110
- Klavžar S, Yero IG (2019) The general position problem and strong resolving graphs. Open Math 17:1126–
- Klavžar S, Rus G (2021) The general position number of integer lattices. Appl Math Comput 390:125664 Korže D, Vesel A (2023) General position sets in two families of Cartesian product graphs. Mediterr J Math 20:203
- Kuziak D, Yero IG, Rodríguez-Velázquez JA (2016) Closed formulae for the strong metric dimension of lexicographic product graphs. Discuss Math Graph Theory 36:1051–1064
- Kuziak D, Puertas ML, Rodríguez-Velázquez JA, Yero IG (2018) Strong resolving graphs: the realization and the characterization problems. Discrete Appl Math 236:270–287
- Manuel P, Klavžar S (2018) A general position problem in graph theory. Bull Aust Math Soc 98:177–187 Oellermann OR, Peters-Fransen J (2007) The strong metric dimension of graphs and digraphs. Discrete Appl Math 155:356–364
- Patkós B (2020) On the general position problem on Kneser graphs. Ars Math Contemp 18:273–280
- Sonnemann E, Krafft O (1974) Independence numbers of product graphs. J Comb Theory Ser B 17:133–142 Thomas EJ, Chandran S.V, Di Stefano G (2024) On the general position number of Mycielskian graphs. Discrete Appl Math 353:29–43
- Tian J, Xu K (2021) The general position number of Cartesian products involving a factor with small diameter. Appl Math Comput 403:126206
- Tian J, Klavžar S (2025) Variety of general position problems in graphs. Bull Malays Math Sci Soc 48:5
- Tian J, Xu K, Klavžar S (2021) The general position number of Cartesian product of two trees. Bull Aust Math Soc 14:1–10
- Yao Y, He M, Ji S (2022) On the general position number of two classes of graphs. Open Math 20:1021-1029

Publisher's Note Springer Nature remains neutral with regard to jurisdictional claims in published maps and institutional affiliations.

