

**GRAPHS WITH TOTAL MUTUAL-VISIBILITY NUMBER  
ZERO AND TOTAL MUTUAL-VISIBILITY IN  
CARTESIAN PRODUCTS**

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**Abstract**

If  $G$  is a graph and  $X \subseteq V(G)$ , then  $X$  is a total mutual-visibility set if every pair of vertices  $x$  and  $y$  of  $G$  admits a shortest  $x, y$ -path  $P$  with  $V(P) \cap X \subseteq \{x, y\}$ . The cardinality of a largest total mutual-visibility set of  $G$  is the total mutual-visibility number  $\mu_t(G)$  of  $G$ . Graphs with  $\mu_t(G) = 0$  are characterized as the graphs in which every vertex is the central vertex of a convex  $P_3$ . The total mutual-visibility number of Cartesian products is bounded and several exact results proved. For instance,  $\mu_t(K_n \square K_m) = \max\{n, m\}$  and  $\mu_t(T \square H) = \mu_t(T)\mu_t(H)$ , where  $T$  is a tree and  $H$  an arbitrary graph. It is also demonstrated that  $\mu_t(G \square H)$  can be arbitrary larger than  $\mu_t(G)\mu_t(H)$ .

**Keywords:** mutual-visibility set, total mutual-visibility set, bypass vertex, Cartesian product of graphs, tree.

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## 1. INTRODUCTION

Let  $G = (V(G), E(G))$  be a graph and  $X \subseteq V(G)$ . Then two vertices  $x$  and  $y$  of  $X$  are  $X$ -visible, if there exists a shortest  $x, y$ -path  $P$  such that  $V(P) \cap X = \{x, y\}$ . The set  $X$  is a *mutual-visibility set* if its vertices are pairwise  $X$ -visible. The cardinality of a largest mutual-visibility set is the *mutual-visibility number*  $\mu(G)$  of  $G$  and a largest mutual-visibility set is called a  $\mu$ -set of  $G$ .

These concepts were introduced and studied for the first time by Di Stefano in [4]. The study was motivated in many ways, notably by the role that mutual-visibility plays in problems arising in the context of distributed computing by mobile entities, and by the fact that vertices in mutual-visibility may represent entities on some nodes of a network that want to efficiently communicate in such a way that the messages do not pass through other entities. We also mention related concepts in computer science that have been explored: distributed computing by mobile entities [5], mutual-visibility tasks [3], and fat entities modelled as disks in the Euclidean plane [12]. A related graph theory topic is the general position in graphs, introduced in [10, 15] and extensively studied by now, cf. [11, 17]. The general position problem was investigated in detail on Cartesian product graphs [6, 8, 9, 13, 14].

In [1], the mutual-visibility problem was studied on Cartesian products and on triangle-free graphs, while in [2] the focus was on strong products. In these studies, the following tools have proven to be extremely useful. We say that  $X \subseteq V(G)$  is a *total mutual-visibility set* of  $G$  if every pair of vertices  $x$  and  $y$  of  $G$  is  $X$ -visible, that is, there exists a shortest  $x, y$ -path  $P$  with  $V(P) \cap X \subseteq \{x, y\}$ . Note that, by definition, the empty set is a total mutual-visibility set. The cardinality of a largest total mutual-visibility set of  $G$  is the *total mutual-visibility number*  $\mu_t(G)$  of  $G$ . Hence, if the empty set is the only total mutual-visibility set of  $G$ , then  $\mu_t(G) = 0$ . Further,  $X$  is a  $\mu_t$ -set if it is a total mutual-visibility set with  $|X| = \mu_t(G)$ .

As observed in [1], there exist graphs  $G$  with  $\mu_t(G) = 0$ . Partial results on such graphs were proved, in particular cactus graphs  $G$  with  $\mu_t(G) = 0$  were characterized. In Section 3 we characterize general graphs  $G$  for which  $\mu_t(G) = 0$  holds as the graphs that contain no bypass vertices. We introduce the latter concept in Section 2, where we also give further definitions needed, recall some known results, and add a few additional preparatory results. In Section 4 we prove bounds on the total mutual-visibility number of Cartesian product graphs and demonstrate their sharpness by several exact results. For instance,  $\mu_t(K_n \square K_m) = \max\{n, m\}$  and  $\mu_t(T \square H) = \mu_t(T)\mu_t(H)$ , where  $T$  is a tree. In Section 5 we continue by the investigation of Cartesian products by considering the estimate  $\mu_t(G \square H) \leq \mu_t(G)\mu_t(H)$ . It holds in many cases, but on the other hand we show that  $\mu_t(G \square H)$  can be arbitrary larger than  $\mu_t(G)\mu_t(H)$ . We con-

clude by suggesting some open problems and directions for further investigation.

2. PRELIMINARIES AND BYPASS VERTICES

We recall here some definitions, for other undefined graph theory concepts, see [7]. All the graphs in this paper are simple and connected, unless stated otherwise. The order of a graph  $G = (V(G), E(G))$  is denoted by  $n(G)$ , the minimum degree of  $G$  is denoted by  $\delta(G)$ , and the subgraph of  $G$  induced by  $S \subseteq V(G)$  is denoted by  $G[S]$ .  $S \subseteq V(G)$  is an *independent set* if  $G[S]$  is an edgeless graph. The cardinality of a largest independent set is the *independence number*  $\alpha(G)$  of  $G$ . A subgraph  $H$  of  $G$  is *isometric* if for each vertices  $x, y \in V(H)$  the distance between them is the same in  $H$  and in  $G$ . Further,  $H$  is *convex* if for each vertices  $x, y \in V(H)$ , all shortest  $x, y$ -paths in  $G$  lie completely in  $H$ . The *girth*  $g(G)$  of a graph  $G$  with a cycle is the length of a shortest cycle of  $G$ . If  $\tau$  and  $\tau'$  are two graph invariants, then we say that  $G$  is a  $(\tau, \tau')$ -graph if  $\tau(G) = \tau'(G)$ .

The Cartesian product  $G \square H$  of two graphs  $G$  and  $H$  has the vertex set  $V(G \square H) = V(G) \times V(H)$ , vertices  $(g, h)$  and  $(g', h')$  are adjacent if either  $gg' \in E(G)$  and  $h = h'$ , or  $g = g'$  and  $hh' \in E(H)$ . Given a vertex  $h \in V(H)$ , the subgraph of  $G \square H$  induced by the set  $\{(g, h) : g \in V(G)\}$  is a  $G$ -layer and is denoted by  $G^h$ .  $H$ -layers  ${}^gH$  are defined analogously. Each  $G$ -layer and each  $H$ -layer is isomorphic to  $G$  and  $H$ , respectively. Moreover, it is also well-known that each layer of a Cartesian product is its convex subgraph, see [7, Lemma 12.3]. We will use this fact later on many times, sometimes implicitly. More generally, a subgraph  $K$  of a Cartesian product  $G \square H$  is convex if and only if the projections of  $K$  on  $G$  and on  $H$  are convex [7, Proposition 13.3].

We next recall some known results on the (total) mutual-visibility number. (Recall that a block graph is a graph in which all blocks are complete.)

**Proposition 1** [4, Corollary 4.3]. *Let  $T$  be a tree and  $L$  the set of its leaves. Then  $L$  is a mutual-visibility set and  $\mu(T) = |L|$ .*

**Proposition 2** [2, Proposition 3.3]. *Block graphs (and hence trees and complete graphs) and graphs containing a universal vertex are all  $(\mu, \mu_t)$ -graphs.*

**Proposition 3** [2, Proposition 3.1]. *Let  $G$  be a graph. If  $V(G) = \bigcup_{i=1}^k V_i$ , where  $G[V_i]$  is a convex subgraph of  $G$  and  $\mu_t(G[V_i]) = 0$  for each  $i \in [k]$ , then  $\mu_t(G) = 0$ .*

Here and later on,  $[k]$  stands for  $\{1, \dots, k\}$ . The following straightforward fact will be used several times later on.

**Proposition 4.** *If  $X$  is a total mutual-visibility set of a graph  $G$  and  $Y \subseteq X$ , then  $Y$  is also a total mutual-visibility set of  $G$ .*

To conclude the preliminaries we introduce the following concept which appears essential in the investigation of the total mutual-visibility concept. We say that a vertex  $u$  of a graph  $G$  is a *bypass vertex* if  $u$  is not the middle vertex of a convex  $P_3$  in  $G$ . Otherwise,  $u$  is a *non-bypass vertex*. Let  $BP(G)$  denote the set of all bypass vertices of  $G$  and let  $bp(G) = |BP(G)|$ . For instance, if  $n \geq 1$ , then  $BP(K_n) = V(K_n)$  because there are no convex paths  $P_3$  in a complete graph. Hence  $bp(K_n) = n$ . Similarly,  $bp(K_{n,m}) = n + m$  for  $n, m \geq 2$ . Indeed, if  $u, v, w$  induce a  $P_3$  in  $K_{n,m}$ , then since  $n, m \geq 2$ , there exists a common neighbor  $v'$  of  $u$  and  $w$ , where  $v' \neq v$ , hence no  $P_3$  in  $K_{n,m}$  is convex. On the other hand, if  $n \geq 5$ , then  $bp(C_n) = 0$ .

The basic fact on bypass vertices is the following.

**Lemma 5.** *If  $u$  is a non-bypass vertex of a graph  $G$  and  $X$  is a total mutual-visibility set of  $G$ , then  $u \notin X$ .*

**Proof.** Since  $u$  is a non-bypass vertex of  $G$ , it is the central vertex of a convex  $P_3$ . If  $x$  and  $y$  are the neighbors of  $u$  on this  $P_3$  and  $u$  would lie in  $X$ , then  $x$  and  $y$  would not be  $X$ -visible. Hence  $u \notin X$ . ■

Lemma 5 implies that

$$(1) \quad \mu_t(G) \leq bp(G).$$

This bound is sharp. If  $T$  is a tree with  $n(T) \geq 3$ , then using Proposition 1 and Proposition 2 we get  $\mu_t(T) = bp(T)$ . Similarly,  $\mu_t(K_n) = bp(K_n) = n$ . On the other hand, consider complete bipartite graphs  $K_{n,m}$ ,  $n, m \geq 3$ . From [2, Corollary 3.6] and [4, Theorem 4.9] we know that  $\mu_t(K_{n,m}) = \mu(K_{n,m}) = n + m - 2$ , but  $bp(K_{n,m}) = n + m$ . The graph  $G$  from Figure 1 is another sporadic example for which the bound (1) is strict. We have  $\mu_t(G) = 1$  and  $bp(G) = 2$ , where  $BP(G) = \{g_6, g_7\}$ .

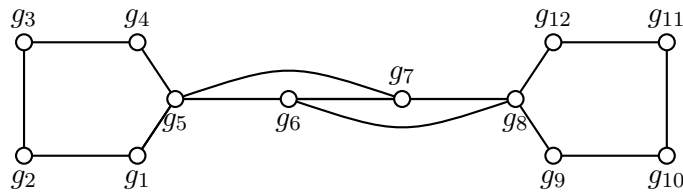


Figure 1. Graph  $G$ .

### 3. GRAPHS WITH $\mu_t = 0$

In this section we characterize graphs with  $\mu_t = 0$  and give several applications of the characterization. We begin by two lemmas, where the first one follows immediately from Proposition 4, and the second one being of independent interest.

**Lemma 6.** *Let  $G$  be a graph. Then  $\mu_t(G) = 0$  if and only if for every  $x \in V(G)$ , the set  $\{x\}$  is not a total mutual-visibility set of  $G$ .*

**Lemma 7.** *Let  $G$  be a graph with  $n(G) \geq 2$  and  $u \in V(G)$ . Then  $\{u\}$  is a total mutual-visibility set of  $G$  if and only if  $u$  is a bypass vertex.*

**Proof.** First assume that  $u$  is a bypass vertex of  $G$ . Then by definition,  $G - u$  is an isometric subgraph of  $G$ . It follows that in  $G$  each pair of vertices is connected by shortest path avoiding  $u$ , hence  $\{u\}$  is a total mutual-visibility set.

The other direction follows by Lemma 5. ■

Note that Lemma 7 does not extend to two vertices. For instance, two opposite vertices of  $C_4$  are bypass vertices, but they do not form a total mutual-visibility set.

The announced characterization now reads as follows.

**Theorem 8.** *Let  $G$  be a graph with  $n(G) \geq 2$ . Then  $\mu_t(G) = 0$  if and only if  $bp(G) = 0$ .*

**Proof.** If  $\mu_t(G) = 0$ , then  $bp(G) = 0$  by Lemmas 6 and 7. Conversely, if  $bp(G) = 0$ , then Lemma 7 says that  $G$  has no singleton total mutual-visibility set, and so by Lemma 6,  $\mu_t(G) = 0$ . ■

Clearly, to check whether a vertex is a bypass vertex is algorithmically simple. Hence the characterization of graphs  $G$  with  $\mu_t(G) = 0$  from Theorem 8 is efficient.

Note that Theorem 8 implies that if  $\mu_t(G) = 0$ , then  $\delta(G) \geq 2$ . Another consequence of the theorem is the following.

**Corollary 9.** *Let  $G$  be a graph with  $g(G) \geq 5$ . Then  $\mu_t(G) = 0$  if and only if  $\delta(G) \geq 2$ .*

The Petersen graph applies to Corollary 9. In addition, the corollary implies the characterization of cactus graphs  $G$  with  $\mu_t(G) = 0$  as given in [2, Proposition 3.2]. For a sporadic example of a graph  $G$  with  $\mu_t(G) = 0$  see Figure 2.

As another application of Theorem 8 we next determine the theta graphs with  $\mu_t = 0$ . For any positive integer  $k \geq 2$  and  $1 \leq p_1 \leq \dots \leq p_k$ , the *theta graph*  $\Theta(p_1, \dots, p_k)$  is the graph consisting of two vertices  $a$  and  $b$  which are joined by  $k$  internally disjoint paths of respective lengths  $p_1, \dots, p_k$ , where  $p_2 \geq 2$ . (We add that several authors use the name theta graph restricted to the case  $k = 3$  in our definition, cf. [19].)

**Corollary 10.** *If  $1 \leq p_1 \leq \dots \leq p_k$ ,  $k \geq 2$ ,  $p_2 \geq 2$ , then  $\mu_t(\Theta(p_1, \dots, p_k)) = 0$  if and only if the following cases hold:*

- (i)  $p_1 = 1$  and  $p_2 \geq 4$ ;

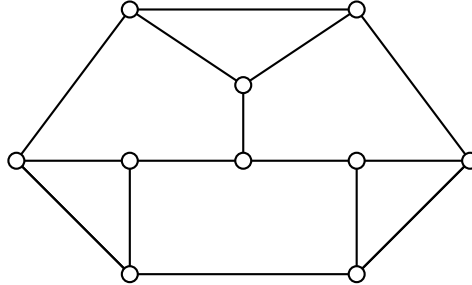


Figure 2. A graph with the total mutual-visibility number 0.

- (ii)  $p_1 = 2$  and  $p_2 \geq 3$ ;
- (iii)  $p_1 \geq 3$ .

**Proof.** Set  $\Theta = \Theta(p_1, \dots, p_k)$  and let  $P_1, \dots, P_k$  be the respective paths of  $\Theta$  connecting  $a$  and  $b$ . Then  $C^i = P_1 \cup P_i$  is an isometric cycle of  $\Theta$  for each  $i \in \{2, \dots, k\}$ . From this fact we infer that  $\mu_t(\Theta) = 0$  if and only if each of the cycles  $C^i$  is of length at least 5. This condition then yields the cases (i)–(iii). ■

We conclude the section with a description of Cartesian products with  $\mu_t = 0$ .

**Theorem 11.** *If  $G$  and  $H$  are graphs, then  $\mu_t(G \square H) = 0$  if and only if  $\mu_t(G) = 0$  or  $\mu_t(H) = 0$ .*

**Proof.** Assume first that  $\mu_t(G \square H) = 0$ . Suppose on the contrary that  $\mu_t(G) \geq 1$  with a total mutual-visibility set  $X$  and  $\mu_t(H) \geq 1$  with a total mutual-visibility set  $Y$ . By Proposition 4, there exist two vertices  $x \in X$  and  $y \in Y$  such that the sets  $\{x\}$  and  $\{y\}$  are total mutual-visibility set of  $G$  and  $H$ . Then we claim that  $U = \{u\}$  with  $u = (x, y)$  is a total mutual-visibility set of  $G \square H$ .

Let  $v, w$  be arbitrary vertices of  $G \square H$ . We need to show that they are  $U$ -visible. If  $v = u$  or  $w = u$ , there is nothing to be proved. If  $v$  and  $w$  lie in the same  $G$ -layer, then  $v$  and  $w$  are  $U$ -visible because their projections onto  $G$  are  $\{x\}$ -visible and since layers in Cartesian product graphs are convex. By the same argument we see that  $v$  and  $w$  are  $U$ -visible when  $v$  and  $w$  lie in the same  $H$ -layer. The last case to consider is when  $v$  and  $w$  neither lie in a common  $G$ -layer or a common  $H$ -layer. Then  $w = (g_1, h_1)$  and  $v = (g_2, h_2)$ , where  $g_1 \neq g_2$  and  $h_1 \neq h_2$ . Then it is well known that there exist two internally disjoint  $v, w$ -shortest paths. Since at least one of these two paths does not contain  $u$ , the vertices  $v$  and  $w$  are  $U$ -visible in  $G \square H$  also in this case. Hence we conclude that  $\mu_t(G \square H) \geq 1$ .

To prove the converse, we may assume, without loss of generality, that  $\mu_t(G) = 0$ . Since  $G^h$  is a convex subgraph of  $G \square H$  for any  $h \in V(H)$ , by Proposition 3, we have  $\mu_t(G \square H) = 0$ . By symmetry, the same result also holds if  $\mu_t(H) = 0$ , as desired. ■

Theorem 11 extends to an arbitrary number of factors as follows.

**Corollary 12.** *If  $G = G_1 \square \cdots \square G_k$ , where  $k \geq 2$ , then  $\mu_t(G) = 0$  if and only if  $\mu_t(G_i) = 0$  for at least one  $i \in [k]$ .*

4. TOTAL MUTUAL-VISIBILITY IN CARTESIAN PRODUCTS

In this section we consider the total mutual-visibility number of Cartesian product graphs. In the previous section we have seen that if  $\mu_t(G) = 0$  or  $\mu_t(H) = 0$ , then  $\mu_t(G \square H) = 0$ . Hence we may restrict our attention here to factor graphs with the total mutual-visibility number at least 1.

To give general bounds we need the following concept. The *independent total mutual-visibility number*  $\mu_{it}(G)$  of  $G$  is the cardinality of a largest independent total mutual-visibility set. Setting  $\ell(G)$  to be the number of leaves of  $G$ , it follows from definitions that for any graph  $G$ ,  $\ell(G) \leq \mu_{it}(G) \leq \min\{\mu_t(G), \alpha(G)\}$ . From Propositions 1 and 2 we know that the leaves set  $L$  is a total mutual-visibility set of a tree  $T$  and  $\mu_t(G) = |L|$ . Hence, if  $n(T) \geq 3$ , then  $\mu_t(T) = |L| = \mu_{it}(T)$ . Note in addition that  $\mu_{it}(K_n) = 1$  while  $\mu_t(K_n) = n$ .

**Theorem 13.** *If  $G$  and  $H$  are graphs of order at least 2,  $\mu_t(G) \geq 1$ , and  $\mu_t(H) \geq 1$ , then*

$$\max\{\mu_{it}(H)\mu_t(G), \mu_{it}(G)\mu_t(H)\} \leq \mu_t(G \square H) \leq \min\{\mu_t(G)n(H), \mu_t(H)n(G)\}.$$

**Proof.** Let  $I_G$  be an independent total mutual-visibility set of  $G$  with  $|I_G| = \mu_{it}(G)$ , and let  $X_H$  be a  $\mu_t$ -set of  $H$ . Set  $U = I_G \times X_H$ . We claim that  $U$  is a total mutual-visibility set of  $G \square H$ .

Let  $(g, h)$  and  $(g', h')$  be arbitrary, different vertices of  $G \square H$ . If  $(g, h)$  and  $(g', h')$  lie in the same  $G$ -layer or the same  $H$ -layer, then  $x$  and  $y$  are  $U$ -visible as layers are convex subgraphs of the product. Hence assume in the rest that  $g \neq g'$  and  $h \neq h'$ .

Let  $g = g_0, g_1, \dots, g_k = g'$  be the consecutive vertices of a shortest  $g, g'$ -path in  $G$  whose internal vertices are not in  $I_G$ . Similarly, let  $h = h_0, h_1, \dots, h_\ell = h'$  be the consecutive vertices of a shortest  $h, h'$ -path in  $H$  whose internal vertices are not in  $X_H$ . Assume first that  $k = 1$ , that is,  $gg' \in E(G)$ . If  $g \notin I_G$ , then the path  $(g, h) = (g_0, h_0), (g_0, h_1), \dots, (g_0, h_\ell), (g_1, h_\ell) = (g', h')$  demonstrates that  $(g, h)$  and  $(g', h')$  are  $U$ -visible. If  $g \in I_G$ , then  $g' \notin I_G$  and then the path  $(g, h) = (g_0, h_0), (g_1, h_0), (g_1, h_1), \dots, (g_1, h_\ell) = (g', h')$  demonstrates that  $(g, h)$  and  $(g', h')$  are  $U$ -visible. Assume in the following that  $k \geq 2$ . Consider now the path  $P$  with the consecutive vertices

$$(g, h) = (g_0, h_0), (g_1, h_0), (g_1, h_1), \dots, (g_1, h_\ell), (g_2, h_\ell), \dots, (g_k, h_\ell) = (g', h').$$

The path  $P$  is a shortest  $(g, h), (g', h')$ -path with no internal vertex in  $U$ . Note that in all the above cases it is possible that  $\ell = 1$  which happens if  $hh' \in E(H)$ .

This proves the claim which implies that  $\mu_t(G \square H) \geq |U| = \mu_{it}(G)\mu_t(H)$ . By the commutativity of the Cartesian product,  $\mu_t(G \square H) \geq \mu_{it}(H)\mu_t(G)$  and the lower bound follows.

Let  $X$  be a total mutual-visibility set of  $G \square H$ . Since each  $G$ -layer  $G^h$  is convex in  $G \square H$  we have  $|X \cap V(G^h)| \leq \mu_t(G)$ , hence  $\mu_t(G \square H) \leq \mu_t(G)n(H)$ . Analogously  $\mu_t(G \square H) \leq \mu_t(H)n(G)$ . ■

In the rest of the section we give several exact results on  $\mu_t(G \square H)$  which also demonstrate that the bounds of Theorem 13 can be attained. We begin with the following sharpness result for the lower bound.

**Corollary 14.** *If  $\mu_t(G \square H) = 1$ , then  $\mu_t(G) = 1$  and  $\mu_t(H) = 1$ .*

**Proof.** By Theorem 11 we have  $\mu_t(G) \geq 1$  and  $\mu_t(H) \geq 1$ . Hence by the lower bound of Theorem 13 we conclude that  $\mu_t(G) = 1$  and  $\mu_t(H) = 1$ . ■

The converse of Corollary 14 does not hold. For instance, consider the theta graph  $\Theta(2, 2, 4)$  as presented in Figure 3.

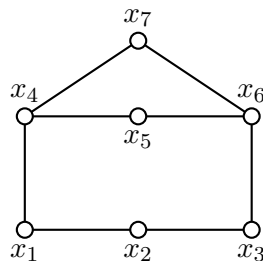


Figure 3. The theta graph  $\Theta(2, 2, 4)$ .

It is straightforward to see that  $\mu_t(\Theta(2, 2, 4)) = 1$  and that  $\{x_5\}$  and  $\{x_7\}$  are  $\mu_t$ -sets. In the product  $\Theta(2, 2, 4) \square \Theta(2, 2, 4)$  one can see that  $\{(x_5, x_5), (x_7, x_7)\}$  is a total mutual-visibility set, hence  $\mu_t(\Theta(2, 2, 4) \square \Theta(2, 2, 4)) \geq 2$ .

[1, Corollary 3.7] asserts that  $\mu(K_n \square K_m) = z(n, m; 2, 2)$ , where  $z(n, m; 2, 2)$  is the Zarankiewicz's number. To determine  $z(n, m; 2, 2)$  is a notorious open problem [16, 18]. Interestingly, the total mutual-visibility number of Cartesian products of complete graphs can be determined as follows, which further demonstrates that the lower bound of Theorem 13 is sharp.

**Proposition 15.** *If  $n, m \geq 2$ , then  $\mu_t(K_n \square K_m) = \max\{n, m\}$ .*



**Proof.** Note first that the vertices of a single  $K_n$ -layer (or a single  $K_m$ -layer) form a total mutual-visibility set. Hence  $\mu_t(K_n \square K_m) \geq \max\{n, m\}$ . To prove the other inequality, set  $V(K_n) = [n]$  and  $V(K_m) = [m]$ , so that  $V(K_n \square K_m) = [n] \times [m]$ . Suppose without loss of generality that  $n \geq m$ . Let  $U$  be an arbitrary total mutual-visibility set of  $G \square H$ . If each  $K_n$ -layer contains at most one vertex of  $U$  there is nothing to be proved. Assume now that some  $K_n$ -layer contains (at least) two vertices of  $U$ . By the symmetry we may assume that  $(1, 1) \in U$  and  $(2, 1) \in U$ . We claim first that  $(i, j) \notin U$ , where  $i, j \geq 2$ . Indeed, if  $(i, j) \in U$ , then the vertices  $(1, j)$  and  $(i, 1)$  are not  $U$ -visible. We claim second that  $(1, j) \notin U$  for  $2 \leq j \leq m$ . Indeed, if  $(1, j) \in U$  for some  $j \geq 2$ , then the vertices  $(1, 1)$  and  $(2, j)$  are not visible. We conclude that if  $(1, 1), (2, 1) \in U$ , then  $U \subseteq V(K_n^1)$  and consequently  $|U| \leq n$ . Analogously we see that if some  $K_m$ -layer contains (at least) two vertices of  $U$ , then  $|U| \leq m$ . In any case,  $\mu_t(K_n \square K_m) \leq \max\{n, m\}$ . ■

The next result (when  $s \in \{3, 4\}$ ) also demonstrates sharpness of the lower bound of Theorem 13.

**Proposition 16.** *If  $s \geq 3$  and  $n \geq 3$ , then*

$$\mu_t(C_s \square K_n) = \begin{cases} 0; & s \geq 5, \\ n; & \text{otherwise.} \end{cases}$$

**Proof.** If  $s \geq 5$ , then Corollary 9 and Theorem 11 yield  $\mu_t(C_s \square K_n) = 0$ . In addition,  $\mu_t(C_3 \square K_n) = n$  by Proposition 15. Hence assume that  $s = 4$  in the remaining proof.

By Theorem 13 we have  $\mu_t(C_4 \square K_n) \geq n$ . It remains to demonstrate that  $\mu_t(C_4 \square K_n) \leq n$ . Let  $R$  be a  $\mu_t$ -set of  $C_4 \square K_n$ . If each  $C_4$ -layer contains at most one vertex of  $R$ , then  $\mu_t(C_4 \square K_n) \leq n$  holds clearly. Suppose next that  $R$  contains at least two vertices from the some  $C_4$ -layer. We may without loss of generality assume that  $(1, n), (2, n) \in R$ . Suppose that there exists another vertex  $(i, j) \in R$ , where  $i \in [4] \setminus [2]$  and  $j \in [n-1]$ . If  $i = 3$ , then the two vertices  $(2, j)$  and  $(3, n)$  are not  $R$ -visible. Similarly, the vertices  $(1, j)$  and  $(4, n)$  are not  $R$ -visible. This would thus mean that  $|R| = 2$ . We conclude that  $\mu_t(C_4 \square K_n) = n$ . ■

**Theorem 17.** *If  $T$  is tree with  $n(T) \geq 3$  and  $H$  is a graph with  $n(H) \geq 2$ , then  $\mu_t(T \square H) = \mu_t(T)\mu_t(H)$ .*

**Proof.** The lower bound  $\mu_t(T \square H) \geq \mu_t(T)\mu_t(H)$  follows by Theorem 13 and the fact that  $\mu_t(T) = \mu_{it}(T)$ .

To prove that  $\mu_t(T \square H) \leq \mu_t(T)\mu_t(H)$ , consider an arbitrary  $\mu_t$ -set  $R$  of  $T \square H$ . Let  $t \in V(T)$  be a vertex with  $\deg_T(t) \geq 2$ . Then  $t$  is a non-bypass vertex of  $T$ . Hence, if  $h \in V(H)$ , then  $(t, h)$  is a non-bypass vertex of  $T \square H$ , thus

$(t, h) \notin R$  by Lemma 5. Therefore,  $R \cap V({}^t H) = \emptyset$ . So  $R$  contains only vertices in  $H$ -layers corresponding to the leaves of  $T$ . Since each such layer can contain at most  $\mu_t(H)$  vertices of  $R$  we conclude that  $\mu_t(T \square H) \leq |R| \leq \mu_t(T)\mu_t(H)$ . ■

As a consequence of Theorem 17 we obtain the following result which demonstrates sharpness of the upper bound of Theorem 13.

**Corollary 18.** *If  $T$  is tree with  $n(T) \geq 3$ , then  $\mu_t(T \square K_n) = n \cdot \mu_t(T)$ .*

### 5. ON THE INEQUALITY $\mu_t(G \square H) \leq \mu_t(G)\mu_t(H)$

All the exact results obtained in Section 4 fulfil the bound

$$(2) \quad \mu_t(G \square H) \leq \mu_t(G)\mu_t(H).$$

Hence one may wonder whether the upper bound of Theorem 13 can be improved/replaced by (2). Before we answer the question, we prove another result where (2) holds.

A graph  $G$  is a *generalized complete graph* if it is obtained by the join of an isolated vertex with a disjoint union of  $k \geq 1$  complete graphs [13]. We further say that  $G$  is a *non-trivial generalized complete graph* if  $k \geq 2$ . Note that if  $G$  is a non-trivial generalized complete graph, then  $\mu_t(G) = n(G) - 1$ .

**Theorem 19.** *If  $G$  and  $H$  are two non-trivial generalized complete graphs, then*

$$\mu_t(G \square H) \leq \mu_t(G)\mu_t(H) = (n(G) - 1)(n(H) - 1).$$

*Moreover, the equality holds if and only if  $G$  or  $H$  is isomorphic to a star.*

**Proof.** Let  $V(G) = \{g_1, \dots, g_{n(G)}\}$  and  $V(H) = \{h_1, \dots, h_{n(H)}\}$ . Let  $g_1$  and  $h_1$  be the universal vertices of  $G$  and  $H$ , respectively.

Note that  $g_1$  is a non-bypass vertex of  $G$  and  $h_1$  is a non-bypass vertex of  $H$ . It follows that each vertex from the layer  ${}^{g_1}H$  is a non-bypass vertex of  $G \square H$  as well as is each vertex from the layer  $G^{h_1}$ . Hence  $G \square H$  contains at least  $n(G) + n(H) - 1$  non-bypass vertices and so by (1),

$$\begin{aligned} \mu_t(G \square H) &\leq bp(G \square H) \\ &\leq n(G)n(H) - (n(G) + n(H) - 1) \\ &= (n(G) - 1)(n(H) - 1). \end{aligned}$$

To prove the equality case, by Theorem 17 we know that  $\mu_t(G \square H) = (n(G) - 1)(n(H) - 1)$  if  $G$  or  $H$  is a star. Suppose in the rest that neither  $G$  nor  $H$  is a star. Then each of them contains an induced subgraph  $K_3$ . Without loss of generality, assume that  $g_1, g_2, g_3$  induce a  $K_3$  of  $G$ , and that  $h_1, h_2, h_3$

induce a  $K_3$  of  $H$ . Then the vertices  $(g_2, h_2), (g_2, h_3), (g_3, h_3)$ , and  $(g_3, h_2)$  induce a  $C_4$  of  $G \square H$ . This  $C_4$  is convex because it is the Cartesian product of  $K_2$  (as a subgraph of  $G$ ) and a  $K_2$  (as a subgraph of  $H$ ) and these two  $K_2$  are clearly convex in respective factors, cf. [7, Proposition 13.3]. Since at most two vertices of this  $C_4$  can lie in a total mutual-visibility set of  $G \square H$ , we conclude that  $\mu_t(G \square H) < bp(G \square H)$ . ■

In the rest we demonstrate that (2) does not hold in general. For this sake we say that a graph  $G$  is *bypass over-visible* if it contains an independent bypass set of vertices  $U$  which contains a  $\mu_t$ -set  $U'$  as a **proper** subset. Note that since  $U'$  is an independent set, a bypass over-visible graph is a  $(\mu_{it}, \mu_t)$ -graph.

Before we state the last result of this paper, we construct two families of bypass over-visible graphs. (Another such family will be presented after the proof of the last result.)

If  $k \geq 3$ , then let  $H_k$  be the graph obtained by attaching two pendant vertices to each of the degree  $k$  vertices of  $K_{2,k}$ . See Figure 4 for  $H_5$ . Let  $U_k$  be the set of all the vertices of  $H_k$  of degree 1 or 2. Then  $U_k$  is an independent bypass set with  $|U_k| = k + 4$ , while every  $\mu_t$ -set of  $H_k$  is obtained from  $U_k$  by removing one degree 2 vertex. Hence each  $H_k$  is a bypass over-visible graphs.

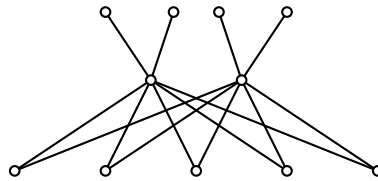


Figure 4. The graph  $H_5$ .

For another family of bypass over-visible graphs let  $\Theta_i$  denote any theta graph  $\Theta(p_1, \dots, p_k)$ , where  $i \geq 2, k \geq 3$ , and

$$2 = p_1 = \dots = p_i < p_{i+1} \leq \dots \leq p_k.$$

Let  $a$  and  $b$  be the vertices of  $\Theta_i$  of degree  $k$ . Then  $BP(\Theta_i)$  consists of the degree 2 vertices which are adjacent to  $a$  and to  $b$ , so that  $bp(\Theta_i) = i$ . Note in addition that  $BP(\Theta_i)$  is an independent set. On the other hand,  $BP(\Theta_i)$  is not a total mutual-visibility set, but becomes such a set if an arbitrary vertex is removed from it. Hence  $\mu_t(\Theta_i) = i - 1$ . We conclude that  $\Theta_i$  is a bypass over-visible graph.

**Theorem 20.** *If  $G$  and  $H$  are bypass over-visible graphs, then*

$$\mu_t(G \square H) > \mu_t(G)\mu_t(H).$$

**Proof.** Since  $G$  and  $H$  are bypass over-visible graphs, there exist independent bypass vertex sets  $I_G$  and  $I_H$  of  $G$  and  $H$ , respectively, which contain  $\mu_t$ -sets  $S_G$  and  $S_H$  as proper subsets. Hence there exist vertices  $u \in I_G \setminus S_G$  and  $v \in I_H \setminus S_H$ . We set  $U = (S_G \times S_H) \cup \{(u, v)\}$  and claim that  $U$  is a total mutual-visibility set of  $G \square H$ .

Consider two arbitrary vertices  $x = (g, h)$  and  $y = (g', h')$  from  $G \square H$ . Suppose first that  $g = g'$ . If  $g \in S_G$ , then  $x$  and  $y$  are  $U$ -visible because  $S_H$  is a total mutual-visibility set of  $H$ . If  $g = u$ , then  $x$  and  $y$  are  $U$ -visible by Lemma 7 applied to  $(u, v)$  and the layer  ${}^g H$ . In all the other cases  $V({}^g H) \cap U = \emptyset$ , hence there is nothing to prove. If  $h = h'$ , then  $x$  and  $y$  are  $U$ -visible by the same argument.

Assume in the rest that  $g \neq g'$  and  $h \neq h'$ . Let  $P_G : g = g_0, g_1, \dots, g_k = g'$  be a shortest  $g, g'$ -path in  $G$  whose internal vertices are not in  $S_G$ . Similar, let  $P_H : h = h_0, h_1, \dots, h_\ell = h'$  be a shortest  $h, h'$ -path in  $H$  whose internal vertices are not in  $S_H$ . The copy of  $P_G$  in the layer  $G^w$  will be denoted by  $P_G^w$  and the copy of  $P_H$  in the layer  ${}^z H$  will be denoted by  ${}^z P_H$ .

Consider first the case  $k = 1$ , that is, when  $gg' \in E(G)$ . Assume first that  $(g_1, h_0) \notin U$ . If  $(u, v) \notin V({}^{g_1} P_H)$ , then the concatenation of the edge  $(g_0, h_0)(g_1, h_0)$  and the path  ${}^{g_1} P_H$  is a required  $x, y$ -path. Suppose next that  $(u, v) \in V({}^{g_1} P_H)$ . Then, because  $g_1 = u \in I_G$  and since  $I_G$  is independent, we infer that  $g_0 \notin I_G$  and thus also  $g_0 \notin S_G$ . Consequently,  $(g_0, h_\ell) \notin U$  and then the concatenation of the path  ${}^{g_0} P_H$  with the edge  $(g_0, h_\ell)(g_1, h_\ell)$  is a required  $x, y$ -path. Assume second that  $(g_1, h_0) \in U$ . Then by the same argument we see that the path  ${}^{g_0} P_H$  followed by the edge  $(g_0, h_\ell)(g_1, h_\ell)$  is again a path which ensures that  $x$  and  $y$  are  $U$ -visible. Similarly we see that  $x$  and  $y$  are  $U$ -visible if  $\ell = 1$ . Note that the argument also applies when  $k = \ell = 1$ .

We are left with the case when  $k \geq 2$  and  $\ell \geq 2$ . Assume first that  $u \neq g_i$  for  $i \in [k - 1]$ . Then the vertices

$$x = (g_0, h_0), (g_1, h_0), (g_1, h_1), \dots, (g_1, h_\ell), (g_2, h_\ell), \dots, (g_k, h_\ell) = y$$

induce a shortest  $x, y$ -path and the internal vertices of it are not in  $U$ . Hence  $x$  and  $y$  are  $U$ -visible. Similarly we see that  $x$  and  $y$  are  $U$ -visible if  $v \neq h_j$  for  $j \in [\ell - 1]$ . The remaining case is that  $u = g_i$  for some  $i \in [k - 1]$  and  $v = h_j$  for some  $j \in [\ell - 1]$ . Then the vertices

$$x = (g_0, h_0), \dots, (g_{i-1}, h_0), (g_{i-1}, h_1), \dots, (g_{i-1}, h_\ell), (g_i, h_\ell), \dots, (g_k, h_\ell) = y$$

induce a shortest  $x, y$ -path in  $G \square H$  with no internal vertices in  $U$ . We conclude that in any case  $x$  and  $y$  are  $U$ -visible. ■

We next present another family of bypass over-visible graphs. Let  $m \geq 1$ . Then we define the graph  $G_m$  as follows. The vertex set is

$$V(G_m) = \{x_0, x_1, \dots, x_{m+2}\} \cup \{y_1, \dots, y_m\} \cup \{z_1, \dots, z_m\}.$$

For any  $i \in [m]$  we connect  $y_i$  and  $z_i$  with  $x_i$  and  $x_{i+1}$ . Finally add the edges  $x_0x_1$  and  $x_{m+1}x_{m+2}$ . See Figure 5 for  $G_5$ .

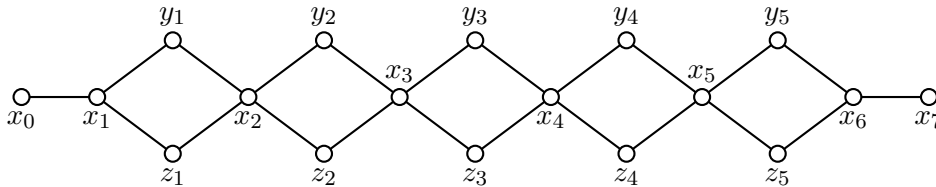


Figure 5. The graph  $G_5$ .

It is straightforward to see that  $BP(G_m) = \{x_0, x_{m+2}\} \cup \{y_1, \dots, y_m\} \cup \{z_1, \dots, z_m\}$ , hence  $bp(G_m) = 2m + 2$  and  $BP(G_m)$  is an independent set. In addition, since for any  $i$ , the vertices  $y_i$  and  $z_i$  cannot both lie in a total mutual-visibility set, the set  $\{x_0, x_{m+2}, y_1, \dots, y_m\}$  is an independent  $\mu_t$ -set of  $G$ . Hence  $G_m$  is a bypass over-visible graph. Moreover,  $bp(G_m) - \mu_t(G_m) = m$ . Now, by a parallel construction as in the proof of Theorem 20 we find out that  $\mu_t(G_m \square G_m) \geq (m + 2)^2 + m$ , hence  $\mu_t(G \square H)$  can be arbitrary larger than  $\mu_t(G)\mu_t(H)$ .

### 6. CONCLUDING REMARKS

There are several possibilities how to continue the investigation of this paper, here we emphasize some of them.

We have characterized the graphs  $G$  with  $\mu_t(G) = 0$ . The next step would be to characterize the graphs  $G$  with  $\mu_t(G) = 1$  (and maybe also with  $\mu_t(G) = 2$ ).

In view of (1) it would be interesting to consider the graphs  $G$  with  $\mu_t(G) = bp(G)$ .

In this paper we had a closer look to the total mutual-visibility number of Cartesian product graphs. Some other graph operations also appear interesting for such investigations, in particular the strong product and the lexicographic product.

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